

# Logarithmic GW th

Steffen Marcus lect 6/17 TMS

## §0 Motivation

$X$  sm proj var.

$GW(X)$ : (virtual) curve counting th.

↳ moduli space of stable maps

$\bar{M}_\Gamma(X)$  proper DM-stack, discrete inv

↳ evaluation

$$\Gamma = g, n, \beta$$

$ev_i: \bar{M}_\Gamma(X) \rightarrow E(X)$  where  $E(X)$  is our evaluation space

↳  $[\bar{M}_\Gamma(X)]^{vir}$  virtual class (usual  $X$ , but for orbifolds or log th get " $E(X)$ ")

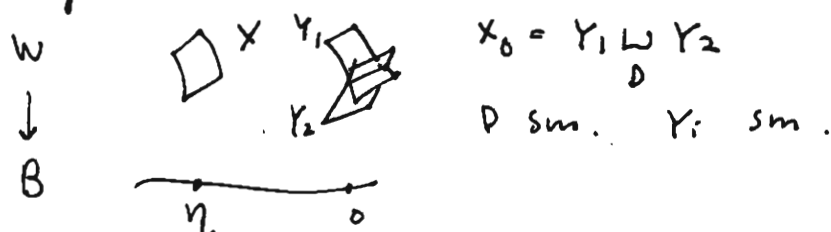
Get:  $\gamma_i \in H^*(E(X))$

$$\langle \gamma_i \rangle_\Gamma = \int_{[\bar{M}_\Gamma(X)]^{vir}} \prod_i ev_i^* \gamma_i$$

counts (virtually) curves in  $X$  w/ incidence conditions with  $[Z_i] = \gamma_i^v$

Techniques: localization, degeneration

Degeneration



Would like:  $GW(X_n) \stackrel{?}{=} GW(X_0)$

Rel  $GW(Y_i/D) \rightarrow$  count (virtually) curves in  $Y_i$  w/ tangency conditions along  $D$ . relative invariants

Regeneration formula:

$$"GW(X) = Rel GW(Y_1/D) * Rel GW(Y_2/D)"$$

some complicate convolution depending on "glueing" along  $D$

P.2 Widely used approach:

"Expanded degenerations"

A-M. Li-Ruan (2001)

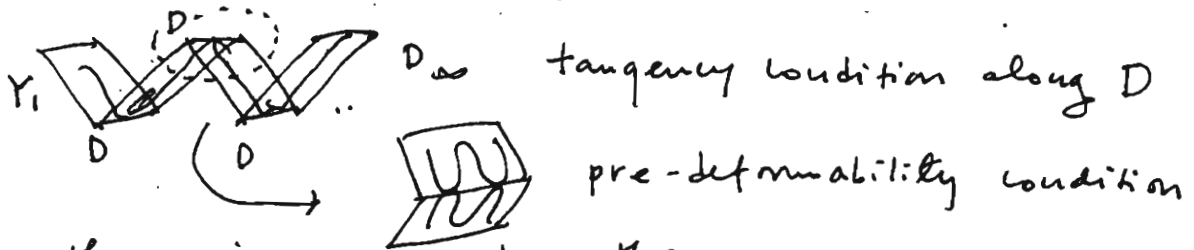
Ionel-Parker (2001/2002)

} Symplectic

Jun Li (2001/2002) Algebraic

Idea deform target  $X$  and come together:

replace  $X_0, Y_i$ 's by "expansions"



The theory is great but there are problems with this th:

- Families of expansions are non-trivial (deformation th is hard)
  - pre-deformability is locally closed
- $D \text{ sm}$  But we want to generalize

Siebert (2001 lecture): Big idea

Use log geometry !!!

Advertisement

Main Existence theorem

(Gross-Siebert 2011, Abramovich-Chen 2011)

2 gps get exactly the same moduli space

$X = (\underline{X}, M_X)$  log variety

log sm  $\underline{X}$  projective

$M_P^{\log}(X) = (\underline{M}_P^{\log}(X), M_{\overline{M}_P(X)})$  stack of log stable map

algebraic,  $\underline{M}_P^{\log}(X)$  is a proper DM stack,

due  $[\underline{M}_P^{\log}(X)]^{\text{vir}}$

### 3 Evaluation map (Abramovich-Chen-Gillan-M) p.3

$$ev: M_{\mathbb{P}^1}^{\log}(X) \longrightarrow \Lambda_{\mathbb{N}}(X)$$

Comparison theorem (Abramovich, -M, -Wise)

$\underline{X}$  sm. prog,  $D$  sm. then

$$Rel. GW(X/D) = Log GW(\underline{X}, M_X)$$

Outline:

- §1 Log geometry
- §2 Log smoothness
- §3 Log curves
- §4 Log stable maps
- §5 Some statements above

In progress: generic formula, degeneration, localization

#### §1 Log Geometry

Monoid: commutative semi-gp w/unit  
 $(+, 0)$  or  $(\cdot, 1)$  (depending on preference)

morphism of monoids  $M \rightarrow N$  st  $0 \mapsto 0$

will define  $X$  log scheme  $(\underline{X}, M_X)$

Let  $\underline{X}$  be a scheme. underlying scheme

Pre-log str Sheaf of monoids  $M_X$  log structure

$M_X$  with a morphism "exponential"

$$M_X \xrightarrow{\alpha} \mathcal{O}_X$$

$(+, 0) \qquad (\cdot, 1)$

Morphisms:  $M_X \longrightarrow N_X$

$$\begin{array}{ccc} & \mathcal{O}_X & \\ \alpha \searrow & \mathcal{O}_X & \swarrow \beta \\ & \mathcal{O}_X & \end{array}$$

Call it a log structure if  $\alpha^{-1}(\mathcal{O}_X^*) \xrightarrow{\sim} \mathcal{O}_X^*$

so  $\mathcal{O}_X^* \subset M_X$

$X = (\underline{X}, M_X)$  is a log scheme.

# 1.4 Characteristic monoid.

$$\bar{M}_X := M_X / \mathcal{O}_X^* \quad (\text{i.e. } 1 \rightarrow \mathcal{O}_X^* \rightarrow M_X \rightarrow \bar{M}_X \rightarrow 0)$$

Why?? eg.  $\begin{array}{ccc} X & \text{sm irred} & \supset D \text{ reduced SNC} \\ \downarrow \text{dominant map} & & \downarrow \\ \Sigma & \text{sm curve} & \supset S \end{array}$   $D = Y_1 \cup \dots \cup Y_m$

Then  $\Omega_X / \Sigma := \Omega_X / f^* \Omega_S$  not locally free where  $f$  is singular

But,  $\Omega_X(\log D)$  diff forms w/ at worst log poles along  $Y_i$ 's

$\uparrow$  then get inclusion  $f^* \Omega_S(\log(S))$  and quotient is locally free!

In terms of "log D",  $f$  is "as good as" smooth

If  $\begin{array}{ccc} m & \xrightarrow{\alpha} & x \neq 0 \text{ in } \mathcal{O}_X \\ \uparrow & & \\ M_X & & \end{array}$  think  $m$  gives a branch at  $\log x$ , will allow " $\frac{dx}{x}$ "

## Associated log structure

given  $\alpha: M_X \rightarrow \mathcal{O}_X$  a pre-log str

define  $(M_X \xrightarrow{\alpha} \mathcal{O}_X)$  as a push out

$$\begin{array}{ccccc} \alpha^{-1} \mathcal{O}_X^* & \hookrightarrow & M_X & \xrightarrow{\alpha} & \mathcal{O}_X \\ & & \downarrow & & \downarrow \\ \mathcal{O}_X^* & \longrightarrow & M_X^a & \longrightarrow & \mathcal{O}_X^* \end{array}$$

$$(m, u) \mapsto \alpha(m) \cdot u$$

$$M_X^a = M_X \oplus \mathcal{O}_X^* / \{ (d^{-1}(p), p) \mid p \in d^{-1}(\mathcal{O}_X^*) \}$$

give the things that need to be glued to get a log-str

Some Fundamental examples:

- $(\mathbb{Z}, \mathcal{O}_X^*)$  trivial log str

- $(\text{Spec } k, k^* \oplus \mathbb{N})$  standard log-point

$$\alpha: k^* \oplus \mathbb{N} \longrightarrow k$$

$$(u, n) \mapsto u \cdot 0^n = \begin{cases} u & n=0 \\ 0 & n \neq 0 \end{cases}$$

Seems silly but is actually very important

Can similarly make  $Q$  log  $p+$

w/ any "sharp" monoid  $Q$  ( $Q^* = \{0,3\}$ )

here 6/18 lect 2 continued)

More meaningful example

$X/k$  scheme  $D \in X$  Divisorial log structure

$$M_{(X,D)}(U) = \{g \in \mathcal{O}_X(U) \mid g \text{ is invertible on } U \setminus D\}$$

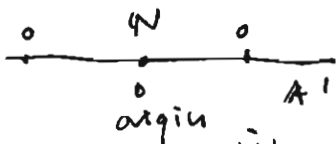
$$\subseteq \mathcal{O}_X(U)$$

$$D = V(x) \subset A^1 = \text{Spec}(k[x])$$

$M_{(A^1,D)}$  generated by  $\mathcal{O}_{A^1}^*$  and  $\{x\}$

$$M_{(A^1,D)} \rightarrow \bar{M}_{(A^1,D)} = i_* N_{D/A^1}$$

$f \mapsto$  order of vanishing at  $x=0$



stalks at  $\bar{M}_{(X,D)} = \begin{cases} N & \text{at } 0 \\ 0 & \text{otherwise} \end{cases}$

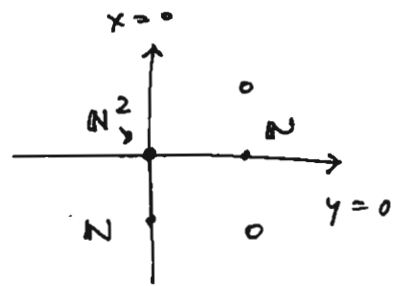
$$D = V(x,y) \subset A^2 = \text{Spec } k[x,y]$$

ii) inclusion of 2 lines resp

$M_{(A^2,D)}$  gen by  $\mathcal{O}_{A^2}^*$  and  $\{x,y\}$

$$M_{(A^2,D)} \rightarrow \bar{M}_{(A^2,D)} = i_* N \oplus j_* N$$

$f \mapsto (\text{ord}_V(x) f, \text{ord}_V(y) f)$



Key example:

$p$  a monoid,  $X$  scheme

$p \rightarrow P(X, \mathcal{O}_X)$  morphism of monoids,  $P_X$  denote the

get  $P_X \xrightarrow{\alpha} \mathcal{O}_X$  a pre-log str. get

"the most sharp on  $X$ "

$\alpha^a : p^a \rightarrow \mathcal{O}_X$  associated log str

$\rightarrow R$  a ring,  $R[p]$  monoid algebra

write  $X = \text{Spec}(D \rightarrow R[p])$

for the underlying scheme,  $X = \text{Spec}(R[p])$

w. asso log  $(p \rightarrow R[p])^a$

p. 6

exercise: describe  $\text{Spec}(\mathbb{N}^r \rightarrow k[\mathbb{N}^r])$ .

To describe the log category, need:

- inverse image log structure given  $f: X \rightarrow Y, M_Y \text{ on } Y$ .  
define  $f^* M_Y$  on  $X$  as  $(f^* M_Y \rightarrow f^* \mathcal{O}_Y \rightarrow \mathcal{O}_X)^{\#}$ .

Now, morphisms of log schemes

a map  $f: X \rightarrow Y$  is  $f: X \rightarrow Y$  and

$$f^{\#}: f^* M_Y \rightarrow M_X \text{ over } \mathcal{O}_X$$

Get LogSch, a category!

eg 1  $D = V(x_1 \cdots x_r) \subset \mathbb{A}^n = X, M_X$  Divisorial log str

$$i: \text{Spec } k \xrightarrow{\text{origin}} X$$

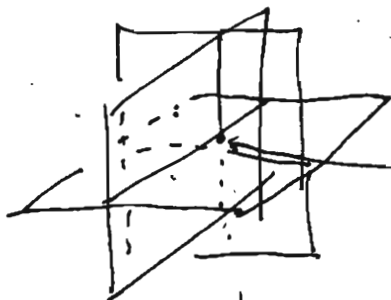
$$i^* M_X = M_X|_0 = \{ \text{germs of } u x_1^{n_1} \cdots x_r^{n_r} \mid n_i \geq 0, u \text{ invertible} \}$$

$$\alpha: i^* M_X \rightarrow i^* \mathcal{O}_X \rightarrow \mathcal{O}_{\text{Spec } k} = k$$

$$u x_1^{n_1} \cdots x_r^{n_r} \mapsto \begin{cases} u(0) & n=0 \\ 0 & n>0 \end{cases}$$

$u = \vec{n}$  (multi-index)

also log str is  $\mathbb{N}^r$  log pt  
at origin  $k^* \oplus \mathbb{N}^r$



Exercise:  $\text{Hom}_{\text{LogSch}}((X, M_X), \text{Spec}(P \rightarrow \mathbb{Z}[P]))$

$\simeq \text{Hom Monoid}(P, \Gamma(X, M_X))$

cont: to last 2 page.

charts a chart for  $P$  on  $X$  is:

$P$ : monoid

$P_X$ : constant sheaf

$$P_X \xrightarrow{\alpha} M_X$$

st.  $(P_X \xrightarrow{\alpha} M_X \rightarrow \mathcal{O}_X)^a$  gives  $P^a \simeq M_X$   $\leftarrow$   
looks a lot like " $M_X \rightarrow \mathcal{O}_X$ " ie

a chart for  $P$  on  $X$

$$P_X \rightarrow M_X$$

$\Downarrow$

$$f: X \rightarrow \text{Spec}(P \rightarrow \mathbb{Z}[P]) \text{ where } f^b: f^* P^a \simeq M_X$$

That is very suitable to think about toric varieties

A chart for  $Q \rightarrow P$  on  $X \rightarrow Y$  is a comm. diagram  
map of monoid

$$\begin{array}{ccc} Q_X & \xrightarrow{\text{chart}} & f^* M_Y \\ \downarrow & \circlearrowleft & \downarrow \text{ on } X \\ P_X & \xrightarrow{\text{chart}} & M_X \end{array}$$

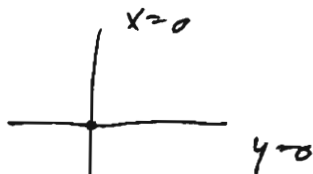
• Caution Here "charts" may not exist at all. It only capture the part of log str where looks like const.

eg node in  $A^2$ ,  $X = V(xy)$ ,  $M_X$  is pull back of

$$\begin{array}{ccc} M(A^2, V(xy)) & & \\ (n,m) \mapsto x^n y^m & & \\ (n,m) \in \mathbb{N}^2 \mapsto k[s,y]/(xy) & & \end{array}$$

$$\begin{array}{ccc} \uparrow & \uparrow \Delta & \uparrow \\ h & \mathbb{N} & k \\ \uparrow & \xrightarrow{\quad} & \uparrow \\ n & \mathbb{N} & \mathbb{N} \end{array}$$

gives a chart on  $f$



$\downarrow f$   
•  $\mathbb{N}$  standard log pt

Exercise: understand this example.

monoid  $p \mapsto p^{\text{gp}} := \{(a,b) \mid (a,b) \sim (c,d) \text{ if } \exists s \in P \text{ st. } s+a+d = s+b+c\}$   
ie. Grothendieck gp of  $P$

integral if  $p \hookrightarrow p^{\text{gp}}$

saturated: integral and  $\forall p \in p^{\text{gp}}, n p \in p \Rightarrow p \in p$   
fine integral & b.g

P.2  
 tonic if fine, saturated, sharp, no torsion  
 fine if  $p \approx \mathbb{N}^r$

A log scheme is "\*" if étale locally  $\exists$  a chart  
 $p \rightarrow M_X$  with  $p$  being with property "\*" .

In the following, assume all monoid are fine, saturated.

Goal: log smoothness

Need:  $\mathcal{L}_X$  log derivations

Follow Grothendieck's formalism.

strict:  $f: X \rightarrow Y$  is strict if  $f^b: f^* M_Y \cong M_X$

eg  $X \rightarrow \mathbb{A}^1, M(X, D) \cong f^* M(\mathbb{A}^1, 0)$

$D = f^{-1}(0) \mapsto 0$ , so  $f$  is strict. \*

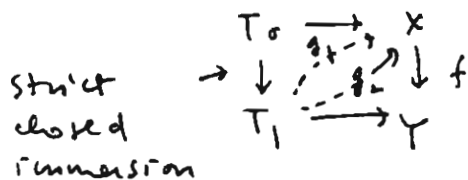
§2. log smoothness

Let  $f: X \rightarrow Y$  be in Log Sch

EGA

Grothendieck/

must consider some lifting problem (eg Hartshorne)



$g_1 - g_2$  a derivation

the problem is "in which category" ?

i.e.  $J$  st  $J^2 = 0$   
 (dual number)

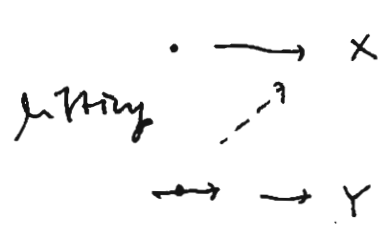
Grothendieck's def<sup>n</sup> of formal smoothness in EGA

say  $T_0 = pt$

(he proved formal sm  $\neq$  sm)

$T_1 = T_0[\epsilon]$

lifting problem on level of



$f^\# = f^* \mathcal{O}_Y \rightarrow \mathcal{O}_X$

$g_1^\# - g_2^\#$  is a scheme derivation

Now, at level of  $f^b, g_1^b - g_2^b$   
 instead of #.



conti on log-derivatives

$$f: X \rightarrow Y$$

log derivation: let  $I$  be an  $\mathcal{O}_X$ -mod

A log-derivation of  $X/Y$  to  $I$  is  $(\partial, D\log)$

•  $\partial: \mathcal{O}_X \rightarrow I \in \text{Der}_Y(X, I)$  ordinary derivation

•  $D\log: M_X \rightarrow I$  additive map st.

$$D\log(ab) = D\log(a) + D\log(b) \quad \forall a, b \in M_X$$

$$\hookrightarrow \alpha(a) D\log(a) = \partial(\alpha(a)) \quad \forall a \in M_X$$

$$\hookrightarrow D\log(a) = 0 \quad \forall a \in f^{-1}(M_Y)$$

ie  $f^* \log f = \text{id}$

log differentials  $\exists$  an  $\mathcal{O}_X$ -module  $\Omega^1_{X/Y}$  with a universal log-derivation  $(\partial, d\log) \in \text{Der}_Y(X, \Omega^1_{X/Y})$

$$\text{st } \forall I, \text{Hom}_{\mathcal{O}_X}(\Omega^1_{X/Y}, I) \xrightarrow{\sim} \text{Der}_Y(X, I)$$

$$u \mapsto (u \circ \partial, u \circ d\log)$$

Explicit construction

$$\Omega^1_{X/Y} = [ \Omega^1_{X/Y} \oplus (\mathcal{O}_X \otimes_{\mathbb{Z}} M_X^{gp}) ] / K$$

where  $K$  is generated by local sections

$$\bullet (d\alpha(a), 0) - (0, \alpha(a) \otimes a) \quad \text{at } M_Y$$

$$\bullet (0, 1 \otimes a) \quad a \in \text{Im}(f^* M_Y \rightarrow M_X)$$

This is a big thm, but will not be proved here.

Remark the 1st one is just  $d \log f = e^{1 \otimes f} d \log f$   
 and just  $D\log(a) = 0$  for  $a$  from the base.

$$\text{Have } \Omega^1_{X/Y}, \quad T_{X/Y} \simeq \mathcal{T}e(\Omega^1_{X/Y}, \mathcal{O}_X)$$

log differential log tangent

Eg  $P$  monoid

$$\bullet X = \text{Spec}(P \rightarrow k[P])$$

$$\text{then } \Omega^1_{X/\text{Spec } k} = \Omega^1_{X/k} \oplus (\mathcal{O}_X \otimes_{\mathbb{Z}} P^{gp}) / K$$

$$\text{Exercise } \nearrow = \mathcal{O}_X \otimes P^{gp}$$

p.2 for  $P = \mathbb{N}^r$ , get  $\Omega_{\mathbb{A}^r/\text{Spec } k}^1 = \mathcal{O}_{\mathbb{A}^r} \otimes \mathbb{Z}^r$

ie.  $\downarrow x_i = x_i (1 \otimes e_i)$

this is the fundamental example

In general,  $h: P \rightarrow Q \rightsquigarrow h^*P: P^{\text{gp}} \rightarrow Q^{\text{gp}}$

$X = \text{Spec}(P \rightarrow \mathbb{Z}[P])$  get  $f: X \rightarrow Y$

$Y = \text{Spec}(Q \rightarrow \mathbb{Z}[Q])$

Then  $\Omega_{X/Y}^1 = \mathcal{O}_X \otimes \text{cok}(h^*P)$

another example

$k[x_1, \dots, x_n] / (x_1, \dots, x_r)$

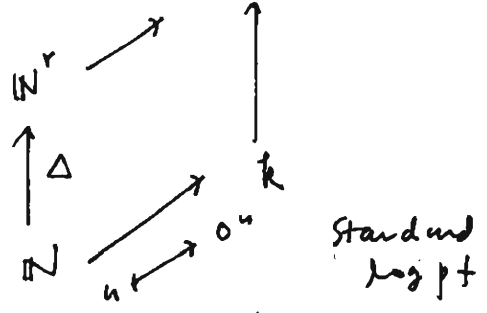
then  $\Omega_{X/Y}^1$ :

generated by

$\frac{dx_1}{x_1}, \frac{dx_2}{x_2}, \dots, \frac{dx_r}{x_r}, dx_{r+1}, \dots, dx_n$

with relation

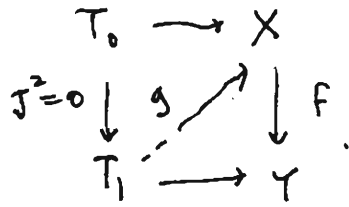
$\sum_{i=1}^r \frac{\downarrow x_i}{x_i} = 0$



Exercise: what if 0

log-smooth:  $f: X \rightarrow Y$  log smooth (log étale)

if  $f$  is loc of finite presentation, and  $\exists$  a sol (nec. unique) to the lifting prob in Log Sch.



Facts: log étale  $\iff$  log sm +  $\Omega_{X/Y}^1 = 0$

log smooth  $\iff$   $\Omega_{X/Y}^1$  locally free

(the  $\Leftarrow$  side is way to be true under more conditions not true even in scheme ??) check

There are 2 Kato's involved in the log th

Kazuya Kato's structure theorems

$f$  log sm  $\Leftrightarrow$  étale locally on  $X$  and  $Y$   
 $f$  fits into  $X \xrightarrow{d_{\text{unt}}} \text{Spec}(\mathbb{Z}[p])$

with  $f \downarrow$   
 $Y \xrightarrow{d_{\text{unt}}} \text{Spec}(\mathbb{Z}[p])$   
 $X \rightarrow Y \times_{\text{Spec}(\mathbb{Z}[p])} \text{Spec}(\mathbb{Z}[p])$  is étale

1) Induced  
 2) kernel and torsion part of  $\text{cok}$  of  $Q^{\text{gp}} \rightarrow p^{\text{gp}}$  are finite grps of order invertible on  $X$

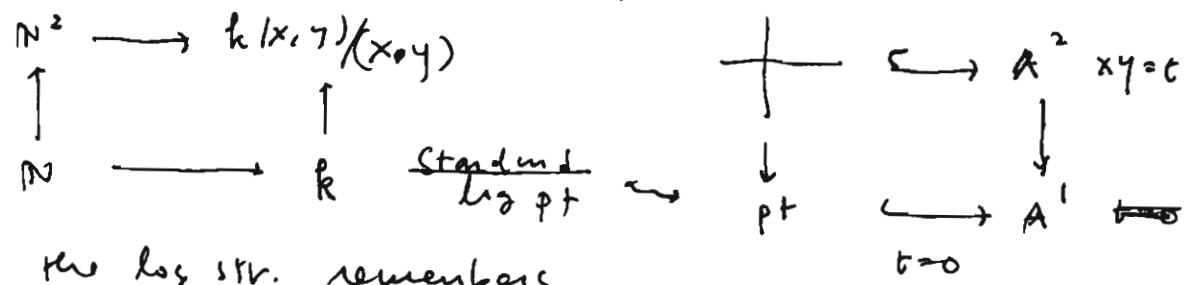
Log sm "means" étale locally toric (toroidal)

eg  $X = \text{Spec}(P \rightarrow k[p])$   
 $Y = \text{Spec} k$  trivial log structure

$X \simeq Y \times_{\text{Spec}(\mathbb{Z}[p])} \text{Spec}(\mathbb{Z}[p])$  and  $X$  is log sm over  $Y$  even if  $X$  is not sm

eg node in  $A^2$ ,  $X = V(xy)$

chart version



i.e. the log str. remembers its deformations (this is the content of the str. thm)

↑  
deform of nodes

Exercise: hot log sm over a trivial log pt

extends thm:  $N \rightarrow P$   
 $1 \rightarrow p$

get  $\text{Spec}(P \rightarrow k[p]) \xrightarrow{f} \text{Spec}(N \rightarrow k[N])$   
 log sm  $X$   $f$   $Y$

Base change  
 weight be  
 reducible and  
 car-reduced

$X_0 \hookrightarrow X$   
 $\downarrow$   
 $\text{Spec} k \hookrightarrow \mathbb{I}$   
 $0: \text{origin}$

gives  $X_0$  log str  $i^* M_X$   
 $\text{Spec} k \cdot 0^* M_Y$   
 get log sm  $X_0 \rightarrow \text{Spec} k$

Warning:  $f$  sm in scheme  $\not\Rightarrow f$  flat  
 but  $f$  log sm in log sch  $\not\Rightarrow f$  flat

examples:  $\mathbb{N}^2 \rightarrow \mathbb{N}^2 \rightarrow \text{Spec } \mathbb{Z}[\mathbb{N}^2] \rightarrow \text{Spec } \mathbb{Z}[\mathbb{N}^2]$   
 $(a,b) \mapsto (a+b, b)$  log sm  
 maybe a not flat: of fine patch of a blow up  
 very ex. check

Integral:  $h: \mathcal{Q} \rightarrow \mathcal{P}$  integral if  $\mathbb{Z}[\mathcal{Q}] \rightarrow \mathbb{Z}[\mathcal{P}]$  is flat, <sup>fine</sup> saturated.

$f: X \rightarrow Y$  integral if  $\forall x \in X, h: (f^{-1}M_y)_x \rightarrow \overline{M_{x,x}}$  is integral

Prop:  $f$  log sm and integral  $\Rightarrow f$  flat.

NB Many semi-stable obj are log sm

Philosophy: MODULI SPACES OF LOG SM OBJECTS (+ INTEGRAL)  
 ARE ALREADY COMPACT !!

(really log sm to imply cpt)

### § 3. Log sm curves

log scheme  
 $\swarrow \searrow$

Def<sup>n</sup>: A log curve is a log sm integral  $f: C \rightarrow W$   
 s.t. geometric fibers are reduced, 1-dim, connected.

Q: What if we don't require fine/satu?

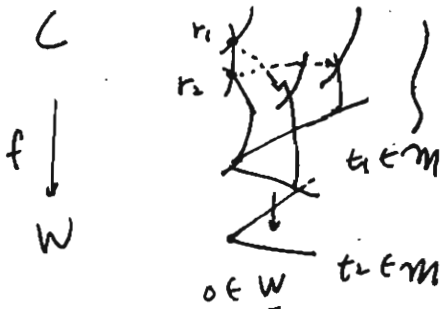
Ex  $P = W \setminus \{1\} = (2,3)$ . Then  $X \cong \text{Spec } k[x,y]/(y^2 - x^3)$   
 $X = \text{Spec}(P \rightarrow k[P])$ .

Furukawa Kato's Structure Theorem (the 2nd Kato)

if you study a log curve  $f: C \rightarrow W$ , get

- singular fibers at worst nodal;
- natural divisorial log str on  $C$  (images of sections)
- natural choices  $p \in M_W$  "smoothing" nodes of  $C_0$  singular

Idea: local description is



partial smoothings

$\underline{W} = \text{Spec } A$   
 $(A, \mathfrak{m})$  complete local ring

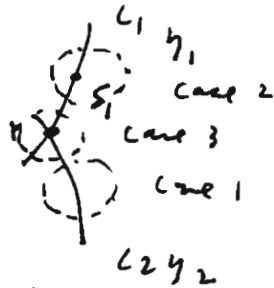
This will give natural log str but the base can have any log str, this cause a problem!

Let  $\varphi = \overline{M}_{W,0}$   $\sigma: \varphi \rightarrow A$  a chart on  $W$

Then étale locally at  $x \in C_0$ ,  $C$  is  $\cong$  to one of  $V \rightarrow W$ .

$x = \eta_i$  1).  $V = \text{Spec } A[z]$  w log str

give pt  $\overline{M}_{C,\eta_i} = \varphi \rightarrow \varphi_V$   
 $\mathfrak{b} \mapsto \sigma(\mathfrak{b})$

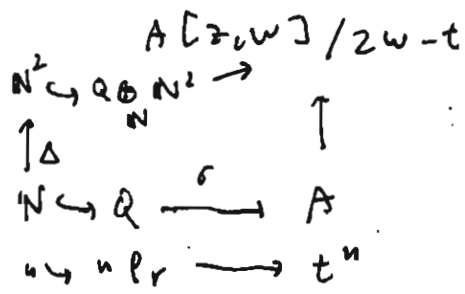


$x = s_j$  2).  $V = \text{Spec } A[z]$

(section)  $\overline{M}_{C,s_j} = \varphi \oplus \mathbb{N} \rightarrow \varphi_V$   
 $(q, a) \mapsto \sigma(\mathfrak{b}) z^a$

$x = r_i$  3).  $V = \text{Spec } A[z, w] / (zw - t)$   $t \in \mathfrak{m}$

$\overline{M}_{C,r_i} = \varphi \oplus \mathbb{N}^2 \rightarrow \varphi_V$  ;  $(\mathfrak{b}, (a, b)) \mapsto \sigma(\mathfrak{b}) z^a w^b$   
 log str on base



$(1,1) = P_r$   
 $z^a = t$   
 are eq'ns at  $\varphi \oplus \mathbb{N}^2$

- get  $P_r \neq 0$  for each node.
- get  $\mathbb{N}$  for section on  $C_0$ .

These 3 charts are essential to do log stable curves.

Notice get "generalization" map

$$\begin{array}{c}
 r \begin{array}{l} \nearrow \eta_1 \\ \searrow \eta_2 \end{array} \\
 \end{array}
 \quad
 \begin{array}{l}
 X_{\eta_i} : \overline{M}_{C,r} \rightarrow \overline{M}_{C,\eta_i} \\
 \mathbb{Q} \oplus \mathbb{N}^2 \xrightarrow{\quad} \mathbb{Q}
 \end{array}$$

$$\eta_1 (a, (a,b)) \mapsto \delta + bPr$$

$$\eta_2 (b, (a,b)) \mapsto \delta + aPr$$

Exercise:  $\iota : \mathbb{Q} \oplus_{\mathbb{N}} \mathbb{N}^2 \xrightarrow{X_{\eta_1}, X_{\eta_2}} \mathbb{Q} \times \mathbb{Q}$  is injective.  
 (good exercise in monoid)

Get a moduli space:

$$M_{g,n}^{\log} = \{ \text{log smooth curve } f : C \rightarrow W \}$$

stable:  $H^0(T_C / \text{spec } k) = 0$  on geom pts

Problem CFG / Log Sch So not yet a stack

Want:  $\overline{M}_{g,n}^{\log} = (\overline{M}_{g,n}^{\log}, M_{\overline{M}_{g,n}^{\log}}^{\log})$  a log alg stack

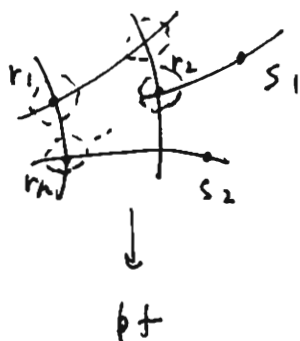
$\rightarrow \overline{M}_{g,n}^{\log}$  DM stack / sch

$M_{\overline{M}_{g,n}^{\log}}^{\log}$  natural / universal

CFG = category  
 fibered in  
 groupoids

For Kato's solution: log curves are nodal, pointed,  
 look a lot like Deligne-Mumford stable.

Build a canonical log str on DM stable curves  $\underline{C}$ .



• Nodes need charts

$$\begin{array}{ccc}
 \mathbb{N}^2 & \longrightarrow & k[z, u, t] / z^4 - t \\
 \Delta \uparrow & & \uparrow \\
 \mathbb{N} & \longrightarrow & k[t]
 \end{array}$$

• div. log str on  $S_i$ 's

$Q = \mathbb{N}^{\# \text{ nodes } (n)}$  on base

$\uparrow$  the smallest number can be chosen

P.4 2 pullbacks :

#1  $f: X \rightarrow S$  a DM stable curve with can log str  $f': X' \rightarrow S'$  a log curve (stable)

Assume

$$\begin{array}{ccc} X' & \xrightarrow{a} & X \\ f' \downarrow & \square & \downarrow f \\ S' & \xrightarrow{b} & S \end{array}$$

then  $a$  and  $b$  extend uniquely to a diagram in log str

$$\begin{array}{ccc} X' & \xrightarrow{a} & X \\ f' \downarrow & \square & \downarrow f \\ S' & \xrightarrow{b} & S \end{array}$$

This is a universal property of "basic" objects in a CFG/log str

call  $f$  satisfying this "basic". Result is

$$\overline{M}_{g,n}^{\log} = \left( \overline{M}_{g,n}^{\text{basic}}, M_{\overline{M}_{g,n}^{\text{basic}}} \right)$$

and  $\overline{M}_{g,n}^{\text{basic}} \cong \overline{M}_{g,n}^{\text{DM}}$  also. Divisorial !!!

So log geom get DM directly without knowing DM's work solves a categorical problem

CFG/log str

Log Al, Stables

$$\{(S, M_S) \rightarrow (X, M_X)\} \longleftarrow (X, M_X)$$

Conversely, given  $y \log$ ,  $\longleftarrow ?$

Answer: study minimality / Basicness

- { F Kato
- Kim (1st attempt to study log curve)
- Gillam

§4. Stable Log Maps :

$X = (X, M_X)$  "nice": log sm  
( $M_X$  SNC Divisorial)

enough for most applications

$$M_{\Gamma}^{\log}(X) = \left\{ \begin{array}{l} \log \text{ curve} \\ \text{(pre-stable)} \end{array} \right. \begin{array}{c} C \xrightarrow{f} X \\ \downarrow \pi \\ U \end{array} \left. \begin{array}{l} \text{underlying maps} \\ \text{of schemes} \\ \text{is stable} \end{array} \right\}$$

would like to understand group points.

On setting:

$k$  alg closed, char  $= 0$

$\underline{W} = \text{Spec } k$

$M_W - Q$  log pt,  $Q$  toric, sharp

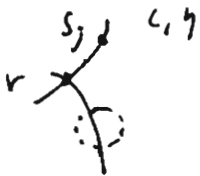
$$k^x \oplus Q \longrightarrow k \quad (u, \gamma) \mapsto \begin{cases} u & \gamma \neq 0 \\ 0 & \gamma = 0 \end{cases}$$

need to understand

$$\begin{array}{ccc} \overline{\pi^* M_W} & \xrightarrow{\psi} & \overline{M_C} \longleftarrow \varphi & \overline{M_X} \\ \uparrow & & & \uparrow \\ \text{stalks at } Q & & & \text{stalks over } P_x \end{array}$$

Let  $x \in C$  be a closed pt

3 possibilities étale locally.



$x$  a smooth pt  $p \in \eta_1, \eta_2$  [1]

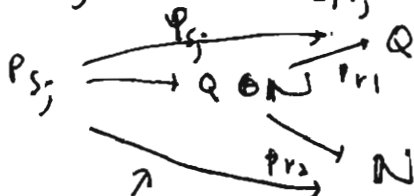
$x$  a marked pt [2]

$x$  is a node [3]

[1]  $x = \eta_i$ . Then  $\psi_{\eta_i}$  is identity and

$$\psi_{\eta_i} : P_{\eta_i} \rightarrow Q \quad \left( \begin{array}{l} \text{map for data on } X \\ \text{for divisor } D, \text{ in } Q \end{array} \right)$$

[2]  $x = S_j$ . Then  $\overline{M_{C, S_j}} = Q \oplus N$ ,  $\psi_{S_j}$  inclusion of  $Q$ ,



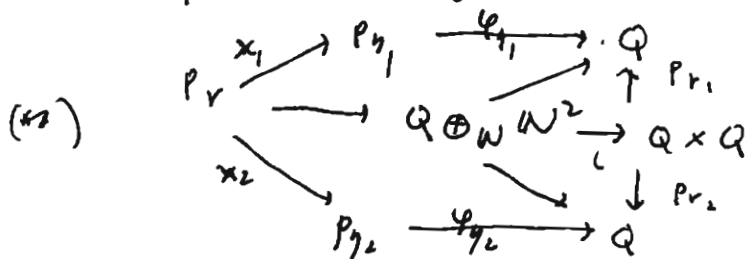
Call this  $\psi_{S_j}$ , then with record contact order data (tangency order)  
(see eg  $X = A^1$  with  $0$  the standard log str) divisors;



P 6

[3]  $x_2 = r$ , then  $\bar{M}_{C,r} = Q \oplus_{\mathbb{N}} \mathbb{N}^2$

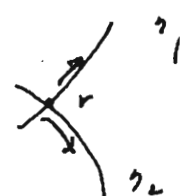
get a diagram



$$\mathbb{N} \xrightarrow{\Delta} \mathbb{N}^2$$

$$\mathbb{N} \rightarrow Q$$

$$1 \mapsto P_r$$



2 generalized maps

Exercise:  $L$  is injective

$$L: Q \oplus_{\mathbb{N}} \mathbb{N}^2 \xrightarrow{x_1 \times x_2} Q \times Q$$

$$(m, (a, b)) \mapsto (m + aP_r, m + bP_r)$$

If  $L(m, (a, b)) = 0$ , then  $m + aP_r = m + bP_r$  with  $P_r \neq 0$

$$(m, (a, b)) = (m, (a, a)) + (m, a(1, -1)) \Rightarrow a = b$$

$$= (m + aP_r, (0, 0)) = 0$$

So get sub manifold

Thus,  $Q \oplus_{\mathbb{N}} \mathbb{N}^2 \hookrightarrow Q \times Q$

in fact, if  $\{ (m_1, m_2) \in Q \times Q \mid m_1, m_2 \in \mathbb{Z}P_r, m \in \mathbb{Q}P_r \}$   
 $m_1 - m_2 = (a-b)P_r$

$$(m_1, m_2) = \begin{cases} L(m_2, (a-b, 0)) & a-b \geq 0 \\ L(m_1, (0, b-a)) & a-b \leq 0 \end{cases}$$

i.e.  $a, b$  are relative "speed" of smoking and wot each comp

get data:  $u_r: P_r \rightarrow \mathbb{Z}$  defined by eq'n

$$\psi_{r_2}(x_2(m)) - \psi_{r_1}(x_1(m)) = u_r(m)P_r \in \mathbb{Z}P_r$$

determined by the commutative diagram (\*)

\*

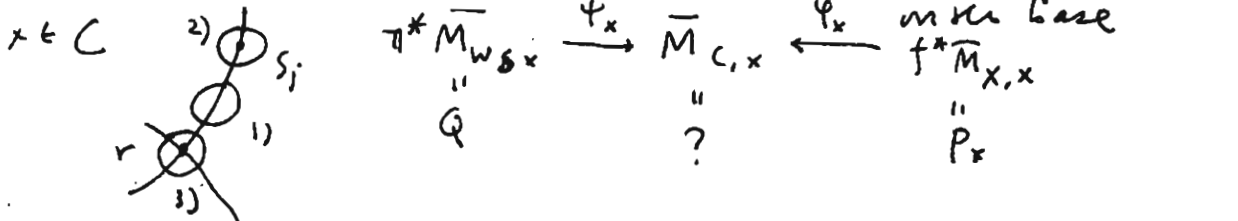
Recall: 6/21 at TMS. 2013

$(X, M_X) = X \text{ log Sm } M_X \text{ -divisorial SNC}$

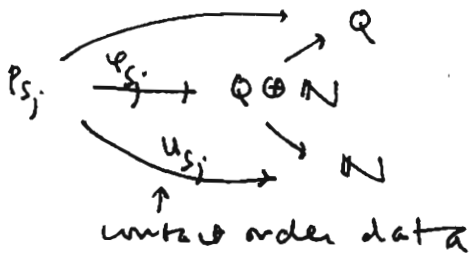
$$M_{\mathbb{P}^2}^{(log)}(X) = \left\{ \begin{array}{l} \text{pre-stable} \\ \text{log curves} \end{array} \downarrow \begin{array}{l} \text{C} \\ \xrightarrow{f} \\ \text{X} \\ \downarrow \pi \\ \text{W} \end{array} \mid \begin{array}{l} \text{Kontsevich stable} \\ \text{of } \mathbb{E}, \mathbb{H} \end{array} \right\}$$

Study geometric pt: (over  $(pt, \mathbb{Q})$ )

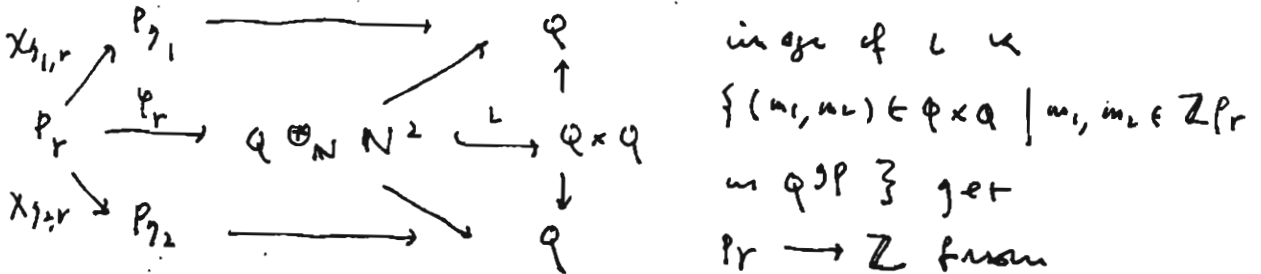
consider



- 1)  $x$  a sm pt  $\eta_i$ , Then  $\bar{M}_{C, \eta_i} = \mathbb{Q}$ ,  $\psi_{\eta_i} = id$ ,  $\psi_{\eta_i}: \mathbb{P}_{\eta_i} \rightarrow \mathbb{Q}$
- 2)  $x$  a marked pt  $S_j$ ,  $\bar{M}_{C, S_j} = \mathbb{Q} \oplus \mathbb{N}$ ,  $\psi_{S_j}$  inclusion of  $\mathbb{Q}$



- 3)  $x$  a node  $r$
- $$\bar{M}_{C, r} = \mathbb{Q} \oplus_{\mathbb{N}} \mathbb{N}^2 \quad \begin{array}{l} (1,1) \\ \mapsto \\ \mathbb{P}_r \end{array}$$



$$\psi_{\eta_2}(x_2(m)) - \psi_{\eta_1}(x_1(m)) = u_r(m) \mathbb{P}_r$$

$\uparrow$   
 $\mathbb{Z}$

ie unit order data of node relative to that component.

Notice Th.3 is the big difference from that of Jun Li's: his pre-definability requires  $r_1=r_2$ , but now could be different



P.2 Def<sup>14</sup>: the type of a stable log map  $f: C \rightarrow X$  is

$$\omega(f: C \rightarrow X) = \left\{ \begin{array}{l} \cdot u_{s_j} \in \text{Hom}(P_{s_j}, \mathbb{N}) \quad \forall j=1, \dots, n \\ \cdot u_{r_k} \in \text{Hom}(P_{r_k}, \mathbb{Z}) \quad \forall \text{ nodes } r_1, \dots, r_k \\ \cdot P_C \text{ dual graph of } C \end{array} \right.$$

this identify all discrete data

Next Q  $\bar{M}_g^{\log}(X) = (\bar{M}_g^{\text{bas}}(X), M_g^{\text{bas}}(X))$   
 what is "basic" log str on base?

Have  $X_{\eta_j, r} P_r \rightarrow P_{\eta_j}$  For  $m \in P_r$ ,

$$a_r(m) := (0 \cdot 0, \underline{X_{\eta_1, r}(m)}, 0 \cdot 0, \underline{-X_{\eta_2, r}(m)}, (0 \cdot 0, u_r(m), 0 \cdot 0))$$

$$\in \left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right)^{\text{gp}}$$

data from target X

power of nodes. data on base

this gives the relation needed

to take care of the data of smoothing a node.

Let  $R \in \left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right)^{\text{gp}}$  saturated subgroup generated by

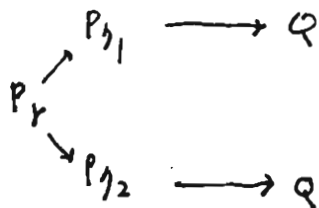
$$\langle a_r(m) \mid r: \text{a node}, m \in P_r \rangle.$$

Def<sup>15</sup>: Basic monoid for a stable log map is:

$$\left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right) \xrightarrow{i} \left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right)^{\text{gp}} \xrightarrow{\pi} \left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right)^{\text{gp}} / R$$

$$Q^{\text{bas}} = [ \text{loc} \left( \prod_{\eta} P_{\eta} \times \prod_r \mathbb{N} \right) ]^{\text{sat}}$$

why??



construct Q using our eq<sup>14</sup>

$$X_{\eta_2, r}(m) - X_{\eta_1, r}(m) = u_r(m) P_r$$

- Before we have Q, there is NO stack can exist!
- This is the most difficult relation in Gross-Siebert program to construct the "basic" Q !!

Have canonical morphisms:

p.3

$$\begin{array}{ccc} \varphi_3 : P_3 & \xrightarrow[\text{morphism of factors}]{} & \prod_3 P_3 \times \prod_1 \mathbb{N} \xrightarrow[\text{"quotient"}]{} Q^{\text{bas}} \\ & & \downarrow \pi \\ & & \mathbb{N} \xrightarrow{} \prod_3 P_3 \times \prod_1 \mathbb{N} \xrightarrow{} Q^{\text{bas}} \end{array}$$

Theorem: This choice satisfies our universal pull back property for "basicness"

Def<sup>n</sup>: A basic s.log Map  $C \xrightarrow{f} X$  is one whose on geom pts of  $W$ ,

The chart on  $W$  is  $Q^{\text{bas}}$ , and  $\varphi, \psi$  are of this type.

(This is Gross-Siebert's version, Qi-Ge Chen also has a independent construction with similar result)

Thm: Basicness is open condition on  $W$

Fact: Basic log str on stable log maps do not introduce new auto morphisms:  $\text{Log Aut } C \subset \text{Aut}$

We have:  $\overline{M}_g^{\text{log}}(x) = (\overline{M}_g^{\text{bas}}(x), \overline{M}_{M_g^{\text{bas}}}(x))$

Need to prove:  $M_g^{\text{bas}}(x)$  alg DM stack.

- Boundedness:  $\left\{ \begin{array}{l} \text{Discrete data are} \\ \beta_i \in H_2(X, \mathbb{Z}), g, n \\ \{ \nu_j : D_j \rightarrow \mathbb{N} \} \end{array} \right.$

fix this, show # of "types" are finite.  
(This is DIFFICULT, in J.Li's module, this is achieved by pro-deformability, here we use "easier" data and get same bdd result)

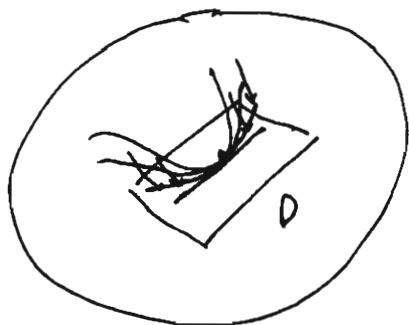
- Properness: Stable reduction (use properness of  $\overline{M}_{g,n}(x, \beta)$ )

properness + locally constant  $\nu_j : P_j \rightarrow \mathbb{N}$

$\Rightarrow$  tangency data remains recorded at limits

- Virtual fund class exists.  
(same theory holds using Behrend Fantechi's construction)

P.4 Picture



a curve falls into the div D

$$v_{s_j} : P s_j \rightarrow Q \oplus \mathbb{N} \rightarrow \mathbb{N}$$

locally constant in families  
in the case, first

$$\begin{array}{ccc} \mathbb{N} & \longrightarrow & \mathbb{N} \\ 1 & \longmapsto & 2 \end{array}$$

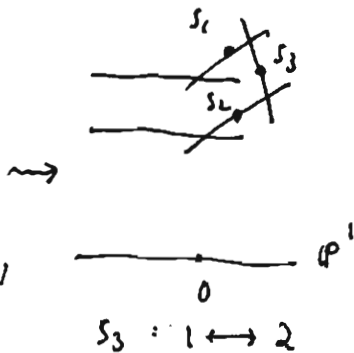
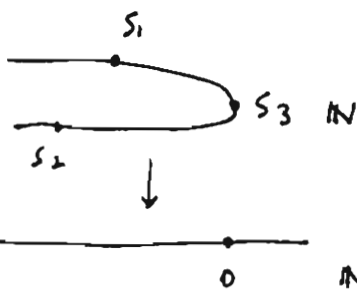
so even when the curve falls into D, the image is still 2

eg.  $X = \mathbb{P}^1 \supset 0$



$$s_3 : \mathbb{N} \rightarrow \mathbb{N} : 1 \mapsto 2$$

$$s_1, s_2 : 0 \rightarrow \mathbb{N} : 0 \mapsto 0$$



J. Li's solution      New solution: fix  
use expanded degenerate by consistency

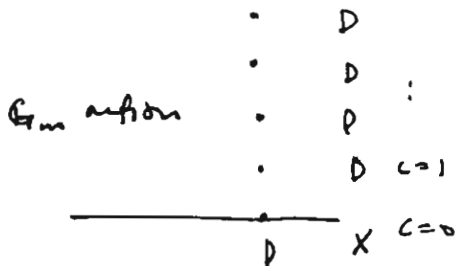
§ 5 Other properties

Evaluation spaces:  $\Lambda_{\mathbb{N}}(X)$

$$ev_i : M_{g,n}^{log}(X) \rightarrow \Lambda_{\mathbb{N}}(X)$$

$$\Lambda_{\mathbb{N}}(X, M_X) = \bigsqcup_{c \in \mathbb{N}} \Lambda_{\mathbb{N}}^c(X)$$

in classical notation



level 0:  $\Lambda_{\mathbb{N}}^0(X) = X \rtimes \mathbb{G}_m$

$[\omega(D)|_0 / \mathbb{G}_m] \leftarrow \Lambda_{\mathbb{N}}^c(X) = \text{"D" twisted by contact order } c.$   
using some (normal) line bundle

$$\beta \mathbb{G}_m = [Pt / \mathbb{G}_m] = \left\{ \begin{array}{l} T \xrightarrow{f} Pt \quad | \quad \mathbb{G}_m \text{ torsor} \\ \downarrow u \quad \quad \quad | \quad f \mathbb{G}_m \text{ equiv} \end{array} \right\}$$

$$= \left\{ \begin{array}{l} L \\ \downarrow u \end{array} \quad | \quad L \text{ a line bundle} \right\}$$

Example:

$$A = [A^1 / \mathbb{G}_m] = \left\{ \begin{array}{l} (L, s) \\ \downarrow u \end{array} \quad | \quad L \text{ line bundle, } s \in P(u, L^{\vee}) \right\}$$

$B\mathbb{G}_m \subset A$  as a divisor at origin

$(A, B, \mathbb{G}_m)$  - universal smooth pair //

Write Relative stable maps

$\overline{Li}(X, D)$  rel maps with / expansions

- expansions
- predeterminability
- difficult obs ths

} 3 hard properties

but there are 4 other theories:

•  $\overline{Kim}(X, D)$  - log maps with expansions

•  $\overline{ACGS}(X, D)$  - log maps

This is what we talk about

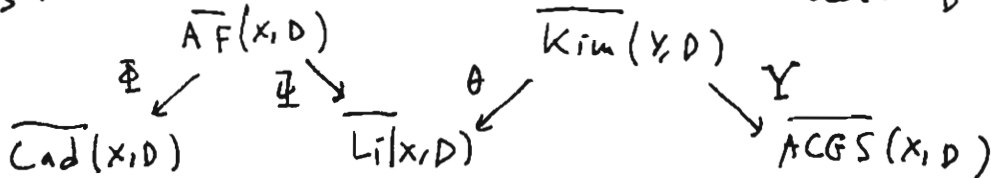
•  $\overline{AF}(X, D)$  - orbifold maps with expansions

•  $\overline{Cad}(X, D)$  - orbifold maps for tangency

All 4 are trying to remove the 3 hard properties

$\overline{ACGS}$  is the 1st to be completely general, namely to allow  $D$  SNC. All others need  $D$  smooth

Maps: let  $(X, D)$  sm pair:



Then (Abramovich - Cadman - Wise 2008 ~)

- $\mathbb{E}_g[AF]^{vir} = [Cad]^{vir} \quad g=0$
- equivalent descendent rel GW(X) in  $g=0$ .

Then (Abramovich, Wise)

- $\mathbb{E}_g[AF]^{vir} = [Li]^{vir}$
- $\mathbb{E}_g[Kim]^{vir} = [Li]^{vir}$  all genus
- $\mathbb{Y}_g[Kim]^{vir} = [ACGS]^{vir}$
- All primary / descenders GW pv's coincide

Diff models, but all the same invariants

1.6. for  $M, M'$  any 2 of these theories,  
Costello's theory (very new)

$$\begin{array}{ccc} M(x, p) & \longrightarrow & M'(x, D) & \text{is Cartesian.} \\ \downarrow & & \downarrow & \\ M(A, B \mathbb{G}_m) & \xrightarrow{\quad \square \quad} & M'(A, B \mathbb{G}_m) & \\ & \text{univ} & & \end{array}$$

gives the comparison of obstruction theory